On regularity and the word problem for free idempotent generated semigroups

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this is what makes people interesting.

And that is the reason we fall in love...

The word problem for semigroups and groups

Definition

A semigroup *S* with a finite generating set *A* has decidable word problem if there is an algorithm which for any two words $w_1, w_2 \in A^+$ decides whether or not they represent the same element of *S*.

Example. $S \cong \langle a, b \mid ab = ba \rangle$ has decidable word problem.

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Some history

- Markov (1947) and Post (1947): first examples of finitely presented semigroups with undecidable word problem;
- Turing (1950): finitely presented cancellative semigroup with undecidable word problem;
- Novikov (1955) and Boone (1958): finitely presented group with undecidable word problem.

The word problem for IG(E)

S - semigroup, E = E(S)

Free idempotent generated semigroup:

$$\mathsf{IG}(E) = \langle E \mid e \cdot f = ef \text{ where } \underbrace{ef = e \text{ or } ef = f \text{ or } fe = e \text{ or } fe = f}_{\text{Basic pairs } (e, f)} \rangle$$

Relates to theory of biordered sets of idempotents (Nambooripad (1979)).

If *E* is finite then this is a finite presentation defining IG(E).

Question Does IG(E) have decidable word problem if *E* is finite? The word problem for IG(E)

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Question

Does IG(E) have decidable word problem if *E* is finite?

Let us consider two illustrative examples:

- 1. Three-element semilattice
- 2. Four-element rectangular band

Example 1: Three-element meet semilattice

S: e f z e f $e^2 = e, f^2 = f, z^2 = z$ ef = fe = ez = ze = fz = zf = z Example 1: Three-element meet semilattice

$$S: \qquad \bigvee_{z}^{e} f \qquad \qquad \underbrace{ \begin{array}{c} \text{Multiplication: } x \cdot y = \inf(x, y) \\ e^{2} = e, \ f^{2} = f, \ z^{2} = z \\ ef = fe = ez = ze = fz = zf = z \end{array}}_{\text{IG}(E) = \langle e, \ f, \ z \mid e^{2} = e, \ f^{2} = f, \ z^{2} = z, \end{array}}$$

$$ez = z, ze = z, fz = z, zf = z \rangle$$

Note: (e, f) not a basic pair $\Rightarrow ef = z$ is not a defining relation.

Example 1: Three-element meet semilattice

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$$ef = fe = ez = ze = fz = zf = z$$
$$\text{IG}(E) = \langle e, f, z \mid e^{2} = e, f^{2} = f, z^{2} = z,$$

$$ez = z, ze = z, fz = z, zf = z\rangle$$

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Example 2: Four-element rectangular band

$$S: \begin{array}{c|c} e_{11} & - e_{12} \\ & \\ e_{21} & - e_{22} \end{array}$$

 $\frac{\text{Multiplication}}{e_{ij}e_{kl} = e_{il}} \quad (\text{note } e_{ij}^2 = e_{ij})$

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► IG(E) \cong {1,2} × \mathbb{Z} × {1,2} where $\mathbb{Z} = \langle a, a^{-1} \rangle$ with multiplication $(i, a^m, j)(k, a^n, l) = (i, a^m p_{jk} a^n, l), \quad P = \begin{pmatrix} 1 & 1 \\ 1 & a \end{pmatrix}.$

This is a Rees matrix semigroup over \mathbb{Z} with structure matrix *P*.

Behaviour exhibited in these examples

Semilattice example

- The regular part of IG(E) is finite and the "same as" the original finite semigroup *S*.
- ► Analysis of the non-regular part of IG(*E*) is necessary to solve the word problem. The word problem is decidable because the non-regular part is well behaved.

Rectangular band example

- There are no non-regular elements in IG(E).
- ▶ Because the subgroups of IG(E) that arise are well behaved (they are all isomorphic to Z and thus have decidable word problem) it follows that IG(E) has decidable word problem.

The general picture





Embedding a group with undecidable word problem

S - semigroup, E = E(S)

Question

Does IG(E) have decidable word problem if *E* is finite?

General facts

- *E* finite \Rightarrow every maximal subgroup of IG(E) is finitely presented.
- ► If IG(*E*) has decidable word problem then every maximal subgroup of IG(*E*) must have decidable word problem.

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Theorem (RG & Ruskuc (2012))

Every finitely presented group is a maximal subgroup of some free idempotent generated semigroup arising from a finite semigroup.

Since there exist finitely presented groups with undecidable word problem...

Corollary

There is a finite semigroup S such that IG(E) has undecidable word problem.

Word problem for regular words

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New question

Does IG(E) have decidable word problem if *E* is finite and every maximal subgroup of IG(E) has decidable word problem?

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Theorem (Dolinka, RG, Ruskuc (2014))

If *E* is finite and every maximal subgroup of IG(E) has decidable word problem then there is an algorithm which, given any two words $u, v \in E^+$

- 1. decides whether both u and v represent regular elements of IG(E) and, if they do,
- 2. decides whether u = v in IG(E).

The general picture





Word problem in general

S - semigroup, E = E(S)

New question

Does IG(E) have decidable word problem if *E* is finite and every maximal subgroup of IG(E) has decidable word problem?

Theorem (Dolinka, RG, Ruskuc (2014))

There exists a finite band $B_{G,H}$ such that:

- (i) All maximal subgroups of $IG(B_{G,H})$ have decidable word problem.
- (ii) The word problem for $IG(B_{G,H})$ is undecidable.

Word problem in general

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For this result we make use of another decision problem...

The membership problem for subgroups

Definition

Let G be a group with finite generating set A, and let H be a subgroup of G given by a finite set of words which generate H.

Then the membership problem for H in G is the problem of deciding, for an arbitrary word w over the generators A, whether or not w represents an element of the subgroup H.

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Theorem (Mihailova (1958))

Their exists a finitely presented group G with a finitely generated subgroup H such that

- ► *G* has decidable word problem, but
- ▶ the membership problem for *H* in *G* is undecidable.

The $B_{G,H}$ construction



Encoding the membership problem



Structure of $IG(B_{G,H})$

Each of \overline{K}'_G and \overline{K}''_G is a Rees matrix semigroup over G

$$\overline{K}'_G \cong I' \times G \times J', \quad \overline{K}''_G \cong I'' \times G \times J''.$$

For any word *w* over *A* the equality $(1', 1, 1')(1'', 1, 1'') = (1', w^{-1}, 1')(1'', w, 1'')$ holds in IG(*B*_{*G*,*H*}) \Leftrightarrow *w* \in *H*.

Conclusion: If $\mathsf{IG}(B_{G,H})$ had decidable word problem this would imply the membership problem for *H* in *G* is decidable, which is a contradiction \pounds

Further questions

S - semigroup, E = E(S) with E finite

Free idempotent generated semigroup:

$$\mathsf{IG}(E) = \langle E \mid e \cdot f = ef \ (e, f \in E, \ \{e, f\} \cap \{ef, fe\} \neq \emptyset) \rangle$$

We have seen that:

- ▶ If all maximal subgroups of IG(E) have decidable word problem then
 - The word problem for regular elements in IG(E) is decidable.
 - The word problem for IG(E) is not decidable in general.

Problem

Find necessary and sufficient conditions for IG(E) to have decidable word problem.

For this it may be useful to investigate the Schützenberger groups of IG(E).